

Category Theory Example Sheet 1

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These questions are of varying difficulty and length. Comments, corrections and clarifications can be emailed to jg352. You can find this sheet on www.dpmms.cam.ac.uk/~jg352/teaching.html.

1. (a) Show that identities in a category are unique.
 (b) Show that a morphism with both a right inverse and a left inverse is an isomorphism.
 (c) Consider $f: A \rightarrow B$ and $g: B \rightarrow C$. Show that if two out of f , g and gf are isomorphisms, then so is the third. [This is known as the *two-out-of-three property*.]
 (d) Show that functors preserve isomorphisms.
 (e) Show that if $F: \mathcal{C} \rightarrow \mathcal{D}$ is full and faithful, and $Ff: FA \rightarrow FB$ is an isomorphism in \mathcal{D} , then $f: A \rightarrow B$ is an isomorphism in \mathcal{C} . [In this case we say F *reflects isomorphisms*.]
2. Let \mathcal{G} be a group viewed as a one-object category. Show that the nat. transformations $\alpha: 1_{\mathcal{G}} \rightarrow 1_{\mathcal{G}}$ correspond to elements in the centre of the group.
3. A morphism $e: A \rightarrow A$ is called *idempotent* if $ee = e$. An idempotent e is said to *split* if it can be factored as fg where gf is an identity morphism.

- (a) Let \mathcal{E} be a class of idempotents in a category \mathcal{C} . Show that there is a category $\mathcal{C}[\check{\mathcal{E}}]$ whose objects are the members of \mathcal{E} , whose morphisms $e \rightarrow d$ are those morphisms $f: \text{dom } e \rightarrow \text{dom } d$ in \mathcal{C} for which $dfe = f$, and whose composition coincides with composition in \mathcal{C} . [Warning: the identity morphism on an object e is not $1_{\text{dom } e}$, in general.]
- (b) If \mathcal{E} is a class of idempotents containing all identity morphisms of \mathcal{C} , show that there is a full and faithful functor $I: \mathcal{C} \rightarrow \mathcal{C}[\check{\mathcal{E}}]$, and that an arbitrary functor $T: \mathcal{C} \rightarrow \mathcal{D}$ can be factored as $\hat{T}I$ for some \hat{T} iff it sends the members of \mathcal{E} to split idempotents in \mathcal{D} .
- (c) If all idempotents in \mathcal{C} split, \mathcal{C} is said to be *Cauchy-complete*; the *Cauchy-completion* $\hat{\mathcal{C}}$ of an arbitrary category \mathcal{C} is defined to be $\mathcal{C}[\check{\mathcal{E}}]$, where \mathcal{E} is the class of all idempotents in \mathcal{C} . Verify that the Cauchy-completion of a category is indeed Cauchy-complete.

4. (a) Show that any functor $F: \mathcal{C} \rightarrow \mathcal{D}$ can be factorised as

$$\mathcal{C} \xrightarrow{L} \mathcal{E} \xrightarrow{R} \mathcal{D}$$

where L is bijective on objects, and R is full and faithful.

- (b) Show that, in a commuting square of functors

$$\begin{array}{ccc} \mathcal{B} & \xrightarrow{F} & \mathcal{D} \\ L \downarrow & & \downarrow R \\ \mathcal{C} & \xrightarrow{G} & \mathcal{E} \end{array}$$

with L bijective on objects and R full and faithful, there exists a unique functor $J: \mathcal{C} \rightarrow \mathcal{D}$ with $JL = F$ and $RJ = G$.

- (c) Deduce that a functor which is both bijective on objects and full and faithful is an isomorphism of categories.
- (d) Deduce that the factorisation in (a) is unique up to unique isomorphism, stating clearly what you take this to mean.

5. Let L be a distributive lattice (i.e. a partially ordered set with finite joins (suprema, \vee) and meets (infima, \wedge), satisfying the distributive law

$$a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$

for all $a, b, c \in L$). Show that there is a category \mathbf{Mat}_L whose objects are the natural numbers, and whose morphisms $n \rightarrow m$ are $m \times n$ matrices with entries from L , where we define ‘multiplication’ of such matrices by analogy with that of matrices over a field, interpreting \wedge as multiplication and \vee as addition. Show also that if L is the two-element lattice $\{0, 1\}$ with $0 \leq 1$, then \mathbf{Mat}_L is equivalent to the category \mathbf{Rel}_f of finite sets and relations between them.

6. Prove that $\theta: \mathbf{Nat}(\mathcal{C}(A, -), F) \rightarrow FA$ from the Yoneda Lemma is natural in F for fixed A .
7. Let \mathcal{C} be a small category, and $F, G: \mathcal{C} \rightarrow \mathbf{Set}$ two functors. Use the Yoneda Lemma to show that a natural transformation $\alpha: F \rightarrow G$ is a monomorphism in $[\mathcal{C}, \mathbf{Set}]$ if and only if all components α_A are monomorphisms in \mathbf{Set} .
8. By an *automorphism* of a category \mathcal{C} , we of course mean a functor $F: \mathcal{C} \rightarrow \mathcal{C}$ with a (2-sided) inverse. We say an automorphism F is *inner* if it is naturally isomorphic to the identity functor. [To see the justification for this name, think about the case when \mathcal{C} is a group.]
- (a) Show that the inner automorphisms of \mathcal{C} form a normal subgroup of the group of all automorphisms of \mathcal{C} . [Don’t worry about whether these groups are sets or proper classes.]
- (b) If F is an automorphism of a category \mathcal{C} with a terminal object 1 , show that $F(1)$ is also a terminal object of \mathcal{C} (and hence isomorphic to 1).
- (c) Deduce that, for any automorphism F of \mathbf{Set} , there is a *unique* natural isomorphism from the identity to F . [Hint: Yoneda]
9. Prove that limits are unique up to unique isomorphism.
10. Consider a commutative diagram of the following form:

$$\begin{array}{ccccc} \cdot & \longrightarrow & \cdot & \longrightarrow & \cdot \\ \downarrow & & \downarrow & & \downarrow \\ \cdot & \longrightarrow & \cdot & \longrightarrow & \cdot \end{array}$$

- (a) (“Pullback composition”) Show that if both small squares are pullbacks, then so is the large rectangle.
- (b) (“Pullback cancellation on the right”) Show that if the large rectangle and the right hand square are pullbacks, then so is the left hand square.
- (c) (“The pullback of a pullback is a pullback.”) Consider a cube in which the back, top and bottom square are pullbacks. Deduce from the above that the front square is also a pullback.
11. A monomorphism $f: A \rightarrow B$ in a category is said to be *strong* if, for every commutative square

$$\begin{array}{ccc} C & \xrightarrow{h} & A \\ g \downarrow & & \downarrow f \\ D & \xrightarrow{k} & B \end{array}$$

with g epic, there exists a (necessarily unique) $t: D \rightarrow A$ such that $ft = k$ and $tg = h$. Show that every regular monomorphism is strong, but that in the finite category represented by the diagram

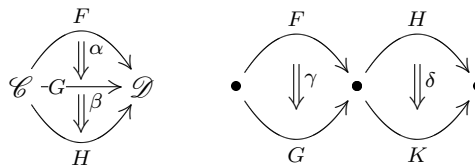
$$\begin{array}{ccc} A & \xrightarrow{f} & B & \xleftarrow{g} & C \\ & \searrow & \downarrow h & \downarrow k & \swarrow m \\ & & D & & \end{array}$$

the morphism f is strong monic but not regular monic.

Additional questions for further practice

These questions will not be discussed in the examples class, they are meant just to provide more opportunity for you to practice and get familiar with the new material.

12. (a) Show that there is a functor $\text{ob}: \text{Cat} \rightarrow \text{Set}$ sending a small category to its set of objects. Is it faithful? Is it full?
- (b) Show that there is a functor $\text{mor}: \text{Cat} \rightarrow \text{Set}$ sending a small category to its set of morphisms. Is it faithful? Is it full?
- (c) Show that the domain and codomain operations give rise to two natural transformations $\text{dom}, \text{cod}: \text{mor} \rightarrow \text{ob}$.
13. In lectures we defined the composite of two natural transformations as on the left below by $(\beta\alpha)_A = \beta_A\alpha_A$. Now consider two natural transformations as on the right below.



Can you work out how to compose these two to get a natural transformation from HF to KG ? Check that your candidate is really natural.

14. Show that the category Set_* of pointed sets is equivalent to the category Part of sets and partial functions.
15. Find representations for the following functors. (All functors are defined on morphisms in the only sensible way.)

- (a) For fixed sets A and B , the functor

$$\begin{aligned} \text{Set}^{\text{op}} &\longrightarrow \text{Set} \\ X &\longmapsto \{\text{pairs of functions } f: X \longrightarrow A \text{ and } g: X \longrightarrow B\}. \end{aligned}$$

- (b) For fixed morphisms $f, g: A \rightarrow B$ in the category Gp , the functor

$$\begin{aligned} \text{Gp}^{\text{op}} &\longrightarrow \text{Set} \\ G &\longmapsto \{\text{morphisms } h: G \longrightarrow A \text{ with } fh = gh\}. \end{aligned}$$

- (c) For a commutative ring R and an ideal I in R , the functor

$$\begin{aligned} \text{CRng} &\longrightarrow \text{Set} \\ S &\longmapsto \{\text{homomorphisms } f: R \longrightarrow S \text{ with } f(I) = 0\}. \end{aligned}$$

16. Recall that a functor from a group \mathcal{G} considered as a one-object category to the category of sets is just a group action. What kind of action does a representable functor $\mathcal{G} \rightarrow \text{Set}$ correspond to?